

# Scalable Data Distribution for Virtual Worlds

## Stanford University

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### Problem:

Today's virtual world systems are characterized by centralized control and limited by the level of scalability they can support. Improved architectures for virtual worlds can enable them to operate more efficiently and scale better to handle millions of users.

### Solution:

Scalability may be achieved by capping the output rate of each virtual world server. We are building a network infrastructure that uses the geometric nature of virtual world interactions to achieve constant bit-rate data traffic from each virtual world server. This approach is physically motivated – humans can perceive a bounded amount of information in real physical spaces, so we could expect similar behavior in a virtual world.

## Constant Bit Rate Servers

### Server Map

Maps sub-regions (spaces) of the virtual world to servers managing them

### OSEG

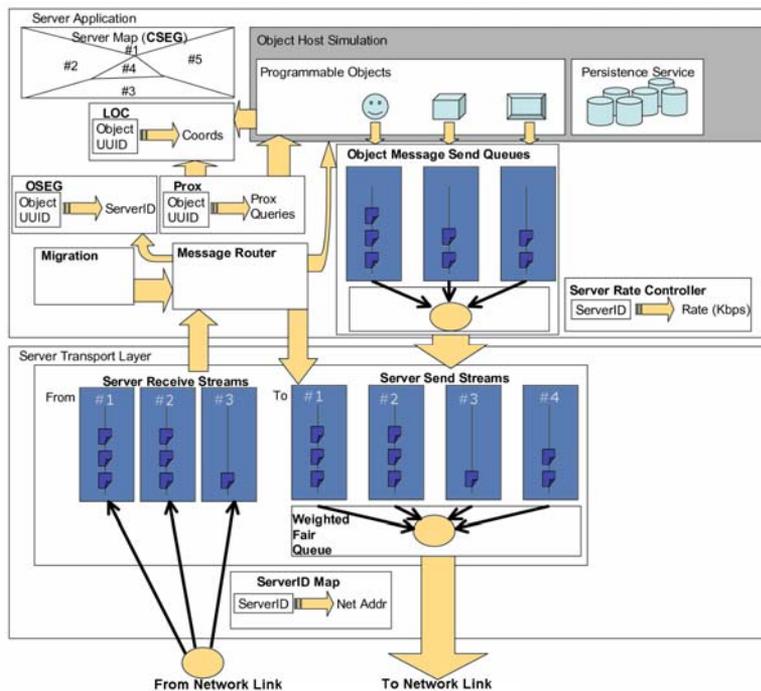
Gets the server on which an object resides

### Migration

Objects migrate when their location changes across regions

### Load balancing

Server map can change dynamically to balance load across servers



### Object Hosts

Objects and clients reside here. In a complete implementation, these would be hosted on remote servers.

### PROX

Introduces nearby objects to each other

### Server Rate Controller

Calculates and controls the rate at which this server can send to another server

### Weighted Fair Queues

Weights proportional to the rates allocated to each server, as given by the Bandwidth Integral

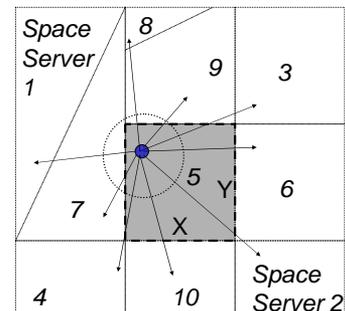
## The Bandwidth Integral

Assume data rate between two objects at distance  $r$  apart =  $f(r)$

$$\text{Total Outgoing Data Rate} = \int_0^X \int_0^Y \int_0^\infty \int_0^{2\pi} f(r) r d\theta dr dy dx \leq \text{constant}$$

Converges if  $f(r) \leq (r^2 \log^{1+\epsilon} r)^{-1}$

$f(r)$  is the data rate allocated for communication to server at distance  $r$



## Performance Metrics

- Location Error – average difference between actual location and sampled location
- Latency of an update through the system
- Data Rate and Fairness of the Weighted Fair Queues

## Work In Progress

- Finish implementation of a full CBR system
- Deploy and evaluate on a large cluster
- Integrate the CBR architecture into the Sirikata virtual world

