

# ON FEEDBACK IN GAUSSIAN NETWORKS

Farzan Farnia, Bobbie Chern, Ayfer Özgür  
 {farnia, bgchern, aozgur}@stanford.edu



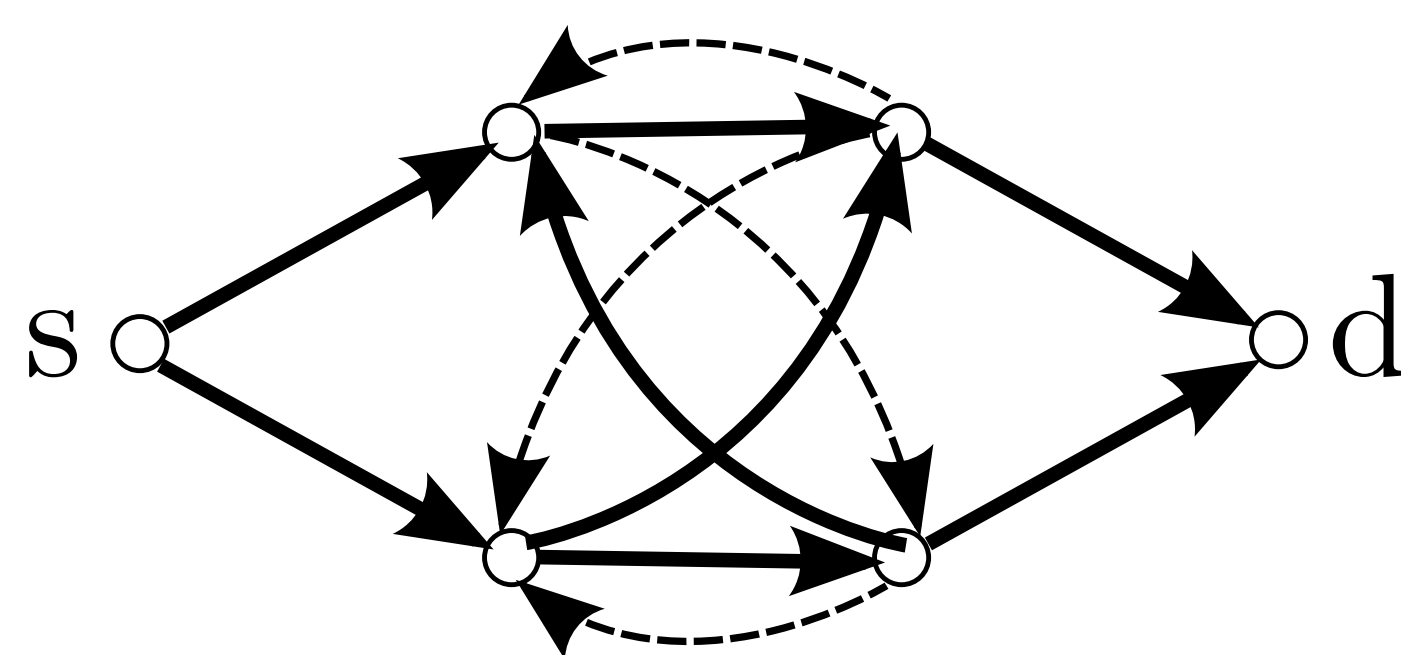
Center for  
 Science of Information  
 NSF Science and Technology Center



## INTRODUCTION

- Feedback has been studied extensively for single-hop communication channels
- There are a lot of opportunities for "feeding back" information in wireless networks
- The nature of these feedback links is significantly different from the feedback models considered in the single-hop settings
- We aim to understand the role of feedback in general Gaussian networks as a model for wireless networks

## MODEL



- We consider a bidirected Gaussian relay network  $G$  consisting of a set of nodes  $V$  and communication links  $E$ .
- We use multiple-input multiple-output channel model, i.e. if we let  $X_v \in \mathbb{C}^{M_v}$  denote the signal transmitted by node  $v \in V$  with  $M_v$  transmit antennas and let  $Y_v \in \mathbb{C}^{N_v}$  denote the signal received by node  $v \in V$  with  $N_v$  receive antennas. We have

$$Y_v = \sum_{u \in V} H_{vu} X_u + Z_v,$$

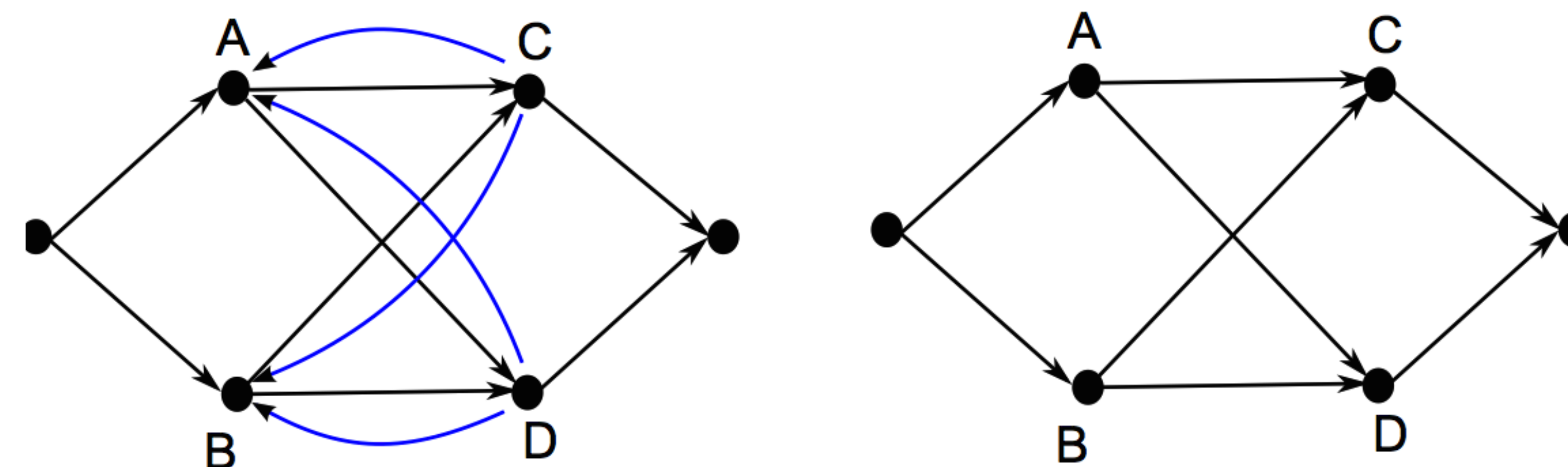
- The noise  $Z_v$  are independent and circularly symmetric Gaussian random vectors  $\mathcal{N}(0, I)$ .
- All nodes are subject to an average power constraint  $P$ .

## TRAFFIC SCENARIOS

We consider the following traffic scenarios over the network:

- Unicast: Source nodes  $s \in V$  wants to communicate to the destination node  $d \in V$ .
- Multiple-Access: Source nodes  $s_1, s_2, \dots, s_n \in V$  want to communicate independent messages to a destination node  $d \in V$ .
- Broadcast: Source node  $s \in V$  wants to communicate independent messages to destination nodes  $d_1, \dots, d_n \in V$ .
- Multicast: Source node  $s \in V$  wants to communicate the same message to destination nodes  $d_1, \dots, d_n \in V$ .
- Multiple-Unicast: Source node  $s_i \in V$  wants to communicate to its destination node  $d_i \in V$  for  $i = 1, \dots, n$ .

## UNICAST



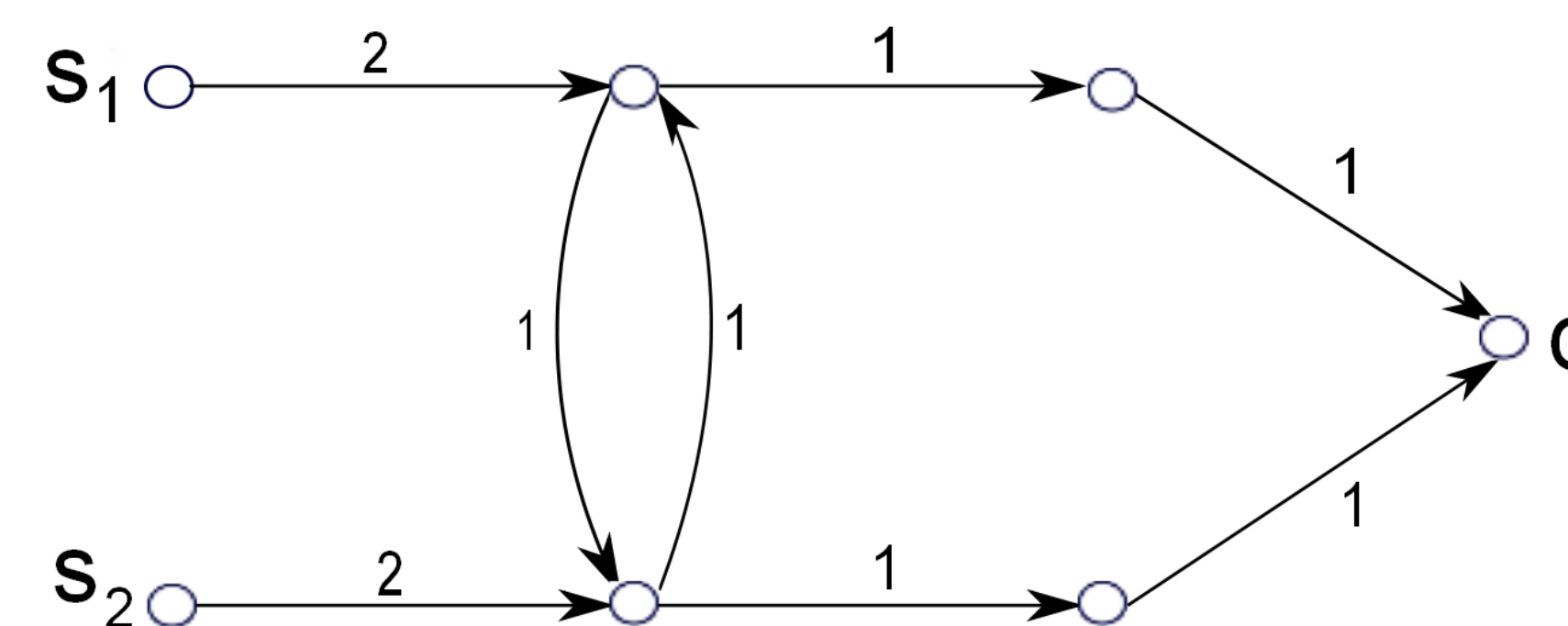
The capacity of the network  $G$ , denoted by  $C(G)$ , is the largest rate at which we can reliably communicate from  $s$  to  $d$ .

**Theorem 1:** In any Gaussian network  $G$  with capacity  $C(G)$ , we can identify a directed acyclic subnetwork  $\tilde{G}$  whose capacity  $C(\tilde{G})$  in bits/s/Hz is bounded by

$$C(G) - g \leq C(\tilde{G}) \leq C(G) + g$$

where  $g = 2 \sum_{v \in V} M_v + (2 + \log 2) \sum_{v \in V} N_v$

## MULTIPLE ACCESS



The capacity region  $C(G)$  is the closure of jointly achievable rate pairs  $R_1, \dots, R_n$  where  $R_i$  is the communication rate from  $s_i$  to  $d$ .

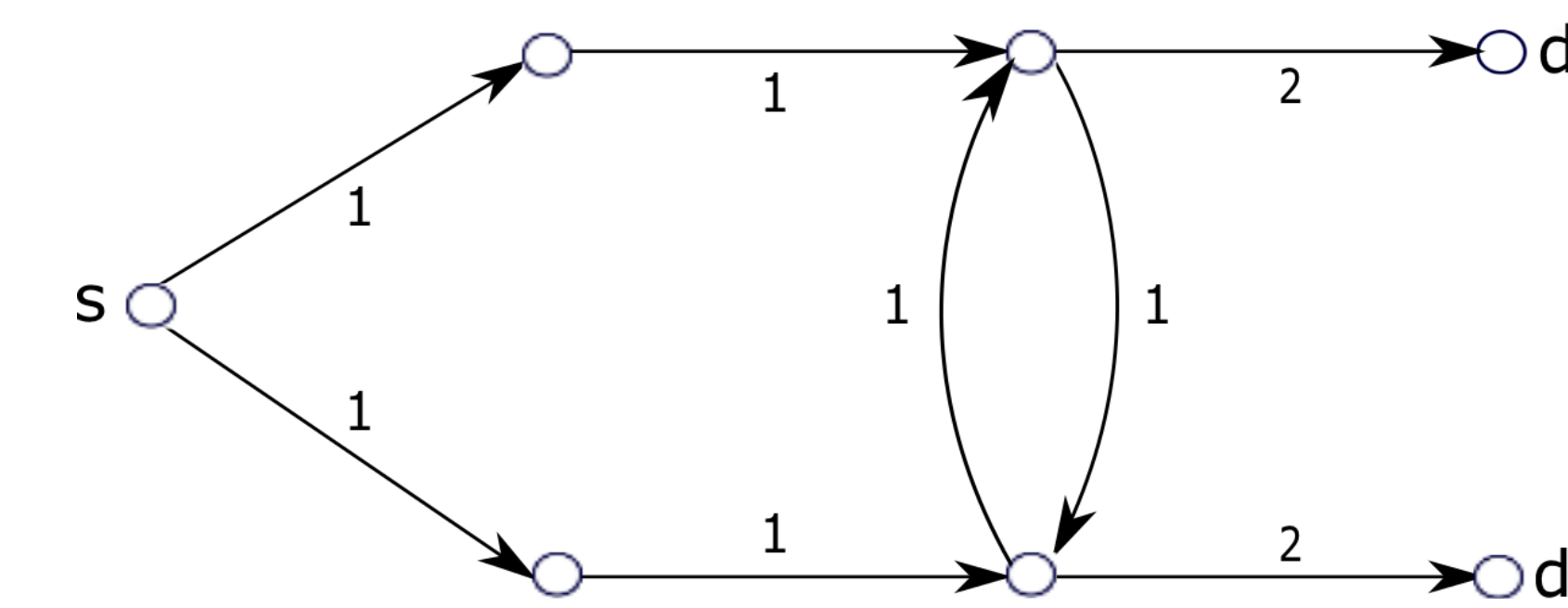
**Theorem 2:** Let  $C(G)$  be the capacity region of a Gaussian network  $G$  with multiple access traffic. If  $(R_1, R_2, \dots, R_n) \in C(G)$ , then there exists an acyclic subnetwork  $\tilde{G}$  such that

$$(R_1 - g_1, R_2 - g_1, \dots, R_n - g_1) \in C(\tilde{G})$$

where  $g_1 = 0.63 \sum_{v \in V} (M_v + N_v)$

However, we cannot conclude that there exists an acyclic subnetwork with the same capacity region as the original network. As a counterexample, see the above figure.

## BROADCAST



The capacity region  $C(G)$  is the closure of jointly achievable rate pairs  $R_1, \dots, R_n$  where  $R_i$  is the communication rate from  $s$  to  $d_i$ .

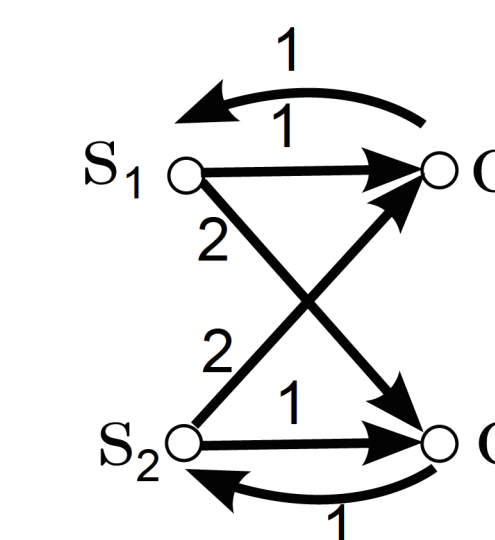
**Theorem 3:** Let  $C(G)$  be the capacity region of a Gaussian network  $G$  with broadcast traffic. If  $(R_1, R_2, \dots, R_n) \in C(G)$ , then there exists an acyclic subnetwork  $\tilde{G}$  such that

$$(R_1 - g_2, R_2 - g_2, \dots, R_n - g_2) \in C(\tilde{G})$$

where  $g_2 = O(M \log M)$ ,  $M = \sum_{v \in V} M_v + \sum_{v \in V} N_v$

However, as in the case of multiple access, the above theorem does not imply the existence of a single acyclic subnetwork whose capacity region is as large as the original network. For a counter example consider the network in the above figure.

## MULTICAST AND MULTIPLE UNICAST



As a counterexample for multicast, consider the network shown for broadcast.

For multiple unicast, the classical Gaussian interference channel with feedback (the above figure) readily provides a counterexample.

## REFERENCES

- 1 Farzan Farnia, Ayfer Özgür, "On feedback in Gaussian multi-hop networks", IEEE Information Theory and Applications Workshop 2014, San Diego
- 2 Bobbie Chern, Ayfer Özgür, "On information flow and feedback in relay networks", IEEE Information Theory Workshop 2013, Seville