A transactional distributed kernel for SERGE

Overview Evaluation

The controller as an OS that safely runs many network applications from many sources (e.g., partially trusted, non-experts)

Model

Programmer sees single-threaded view

Updating hardware = writes to software objects

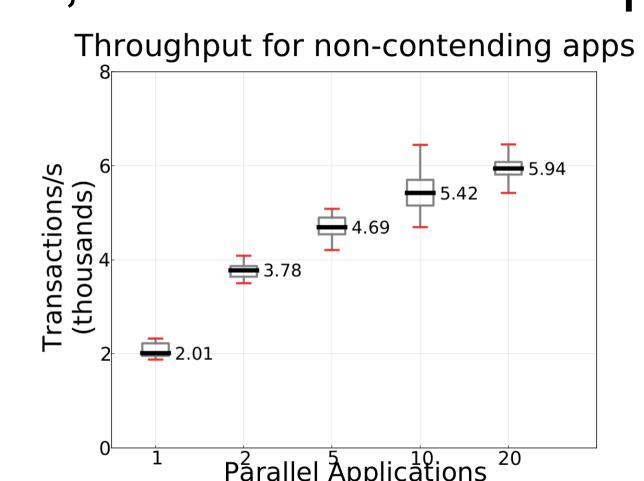
RPCs to other controllers included in transaction (nested, distributed transactions)

Agreement, consistency, isolation, fairness with adequate performance

Single

Latency for contending applications

software defined networks



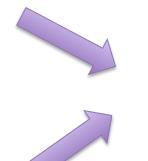
Up to 6k transactions/s on a controller.

Hundreds of transactions/s per switch.



~6.82 ms/program

~4.94 ms/program



Time difference from parallelization of two-phase commit. (Ask about this.)

All experiments in mininet; ~.5ms RTT for multiple controllers

Four Features of SERGEANT

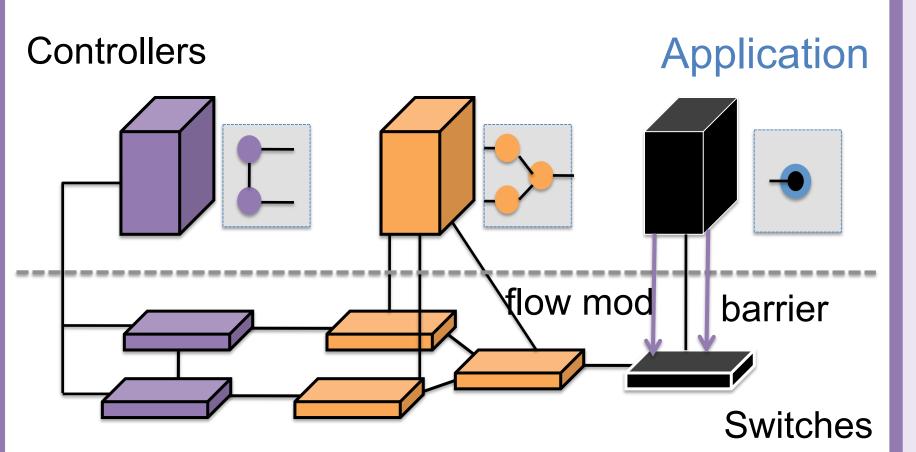
Agreement

between hardware state & software representation of it

Today: Programmer

- Issue barriers
- Handle errors

Sergeant: System-enforced



Automatic transactions to hardware using barriers for dirty state.

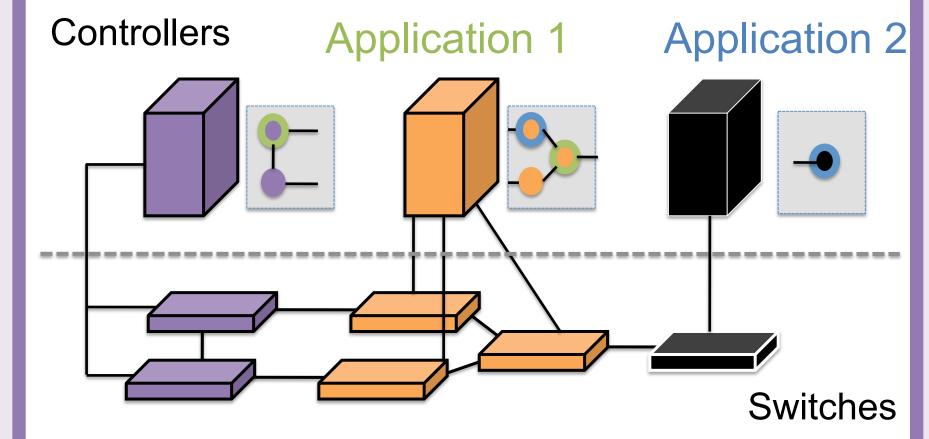
Consistency

of software state as many apps modify it

Today: Programmer

- Locks
- Database

Sergeant: System-enforced



Fine-grained locks on state automatically prevent read/ write conflicts.

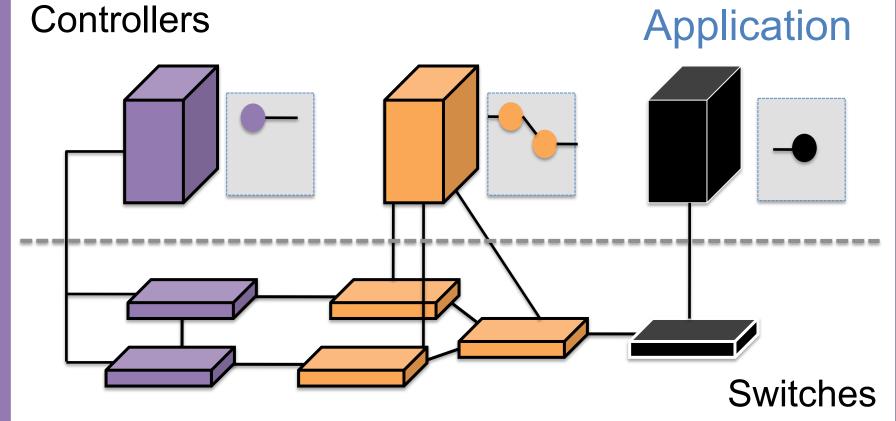
Isolation

between different applications

Today: Programmer

• ???

Sergeant: System-enforced



Permissions hide switches from apps not allowed to operate on them.

Fairness

guaranteed between applications accessing underlying resources

Today: Programmer

• ???

Sergeant: System-enforced

Novel transaction scheduling algorithm provides fairness.

Behram Mistree, Daniel Jackoway, & Philip Levis