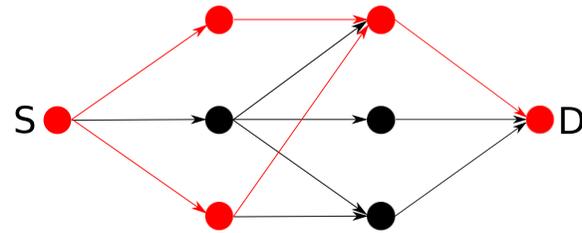




# Wireless Network Simplification

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## Wireless Network Simplification



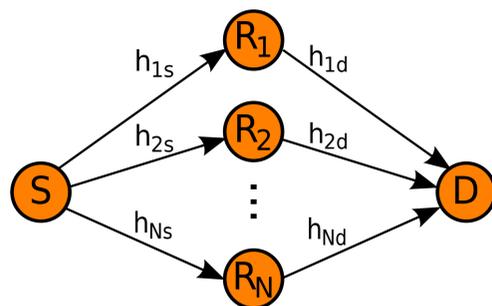
Objective:

Maintain a **good fraction** of the network capacity by using only a **small subset** of the relays.

Motivation

- Using fewer relays allows us to reduce operational complexity.
- Understanding the critical links leads to better resource utilization.
- Unexpectedly, simplification provides better approximations for the capacity of wireless networks than existing results in the literature.

## Gaussian Diamond Network



The source is connected to the relays through a broadcast channel; the relays are connected to the destination through a multiple-access channel with arbitrary channel configurations  $\{h_{ij}\}$ . The source is equipped with  $n_s$  transmit antennas and the destination is equipped with  $n_d$  receive antennas.

## Previous Work

### Single Antenna Diamond Network

*Theorem* [1]: In every every  $N$ -relay diamond network with  $n_s = n_d = 1$ , there exists a subset of  $k$  relays which alone provide a capacity  $C_k$

$$C_k \geq \frac{k}{k+1}C - G,$$

where  $C$  is the capacity of the  $N$ -relay network.  $G$  is a constant independent of the channel configurations and the operating SNR of the network.

By using **only  $k$  relays**, we can achieve approximately a **fraction  $\frac{k}{k+1}$**  of the  $N$ -relay capacity.

### Multiple Antenna Diamond Network

Define  $\alpha_i$  to be the individual capacity from the source to relay  $R_i$  and  $\beta_i$  to be the individual capacity from relay  $R_i$  to the destination.

*Theorem* [2]: When  $\alpha_i = \alpha$  and  $\beta_i = \beta$ ,  $n_s = 2$  and  $n_d = 2$ , in every  $N$ -relay diamond network, there exists a 2-relay subnetwork such that

$$C_2 \geq C - G$$

where  $C$  is the capacity of the entire network and  $G$  is a constant independent of the channel configurations and the operating SNR.

When  $n_s = 2$  and  $n_d = 2$ , we can approximately achieve the capacity of **the whole network** by using **only 2 relays** if **each source-relay relay-destination link has the same capacity**.

The theorem does not hold when  $\alpha_i$  and  $\beta_i$  are arbitrary: we can construct networks with  $n_s = 2$  and  $n_d = 2$  where the largest 2-relay subnetwork only provides half of the capacity of the entire network.

## Main Problem

Extend the simplification framework to wireless networks with **arbitrary** configurations and **arbitrary** number of antennas.

## We Explore

### • MIMO Simplification

Consider a MIMO channel with  $n_t$  transmit and  $n_r$  receive antennas, and capacity  $C_{n_r \times n_t}$ . Assume  $n_r \geq n_t$ . There exists a subset of  $n_t$  receive antennas which achieve a capacity

$$C_{n_t \times n_t} \geq C_{n_r \times n_t} - G$$

where  $G$  is a constant independent of the channel configurations and the operating SNR. An analogous result holds when  $n_t \geq n_r$ .

### • Submodular flows

For any subset of links  $A_i$ , the capacities across those links  $\alpha(A_i)$  satisfies:

$$\begin{aligned} \alpha(A_1) &\leq \alpha(A_2) \quad \text{if } A_1 \subseteq A_2 \\ \alpha(A_1 \cup A_2) &\leq \alpha(A_1) + \alpha(A_2) - \alpha(A_1 \cap A_2) \end{aligned}$$

### • Non-Shannon type inequalities for entropy

Suppose  $X, Y, Z$  are jointly Gaussian random variables. Let  $I(X, Y) = \min(I(X, Y), I(X, Z), I(Y, Z))$ , then:

$$\min(I(X, Z), I(Y, Z)) \leq I(X, Y) + 2$$

Intuitively, if  $X$  and  $Y$  are close to independent, then  $Z$  cannot simultaneously give a lot of information about both  $X$  and  $Y$ .

## References

- [1] C. Nazeroglu, A. Özgür, and C. Fragouli, "Wireless Network Simplification: the Gaussian N-Relay Diamond Network", *IEEE International Symposium on Information Theory (ISIT) 2011*
- [2] C. Nazeroglu, A. Özgür, J. Ebrahimi, and C. Fragouli, "Network Simplification: the Gaussian N-Relay Diamond Network with Multiple Antennas", *IEEE International Symposium on Information Theory (ISIT) 2011*